

Messaging Layer Security



Berlin Crypto Meetup, April 2nd 2019
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Current secure messaging

- **mpOTR, (n+1) sec:** No PCS
- **OpenPGP, S/MIME:** No FS/PCS, linear scaling
- **Pairwise double ratchet:** FS & partial PCS, asynchronous, but linear scaling, no crypto for group membership
(libsignal, Proteus, Olm, etc.)
- **Sender keys:** Linear scaling, PCS possible but very expensive
(WhatsApp, FBM, OMEMO, Megolm, etc.)

Goals and security properties

- **Groups:** support large groups (up to 50k members) efficiently
- **Asynchronous:** No two participants are online at the same time
- **Security:** Manage group membership, modern properties like Forward Secrecy and Post-Compromise Security
- **Formal verification:** similar to TLS, verified specification & implementations
- **Standardized:** Interoperable implementations

History

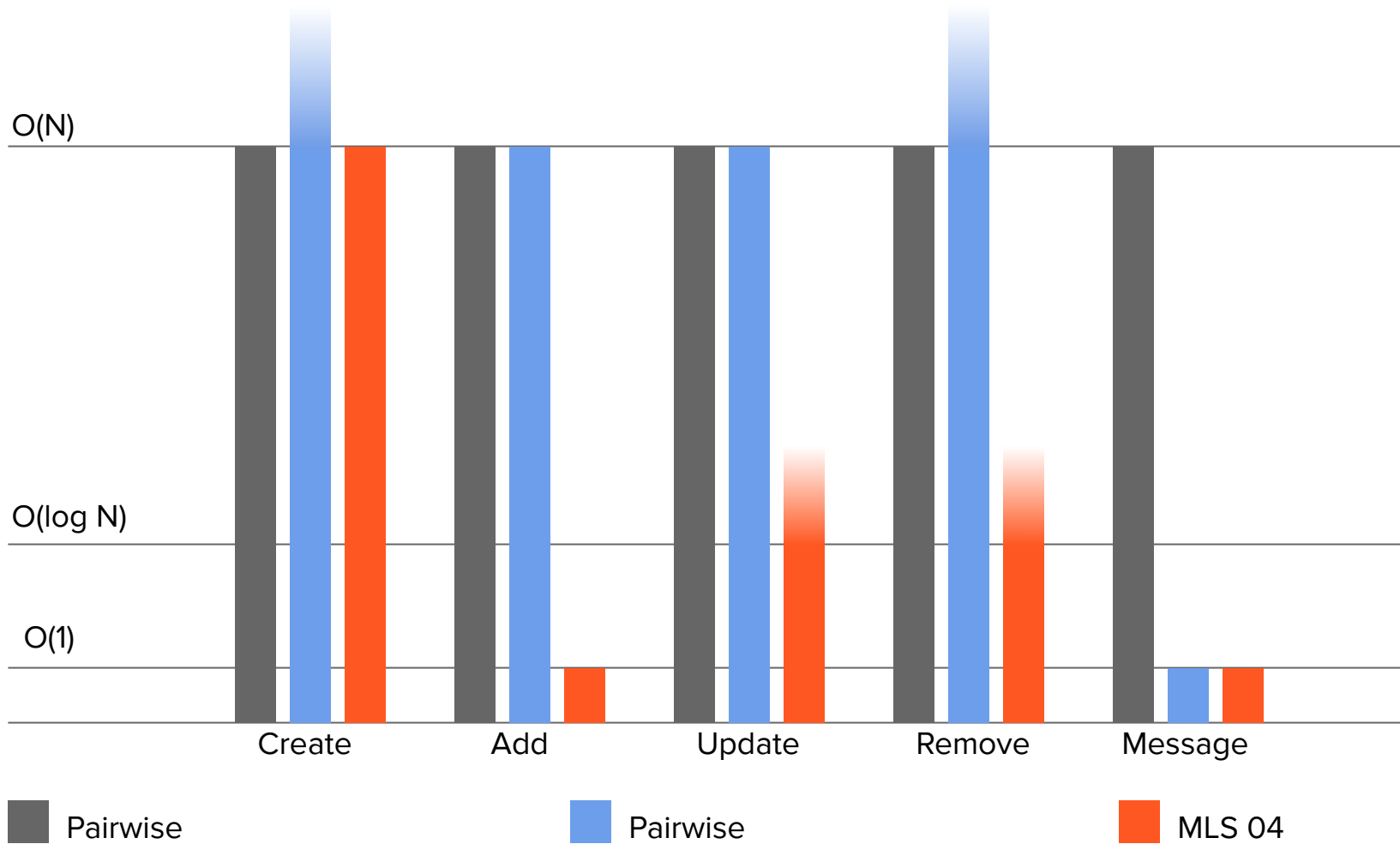
- 2015** Tree-based encryption schemes
(FB, University of Oxford, Mozilla, Cisco, ...)
- 2016** Wire seeks partners for IETF standard
- 2017** Asynchronous Ratcheting Trees ([paper](#))
(Millican, Cohn-Gordon, Cremers, et al)
- 2018** MLS IETF BoF & WG

The IETF working group

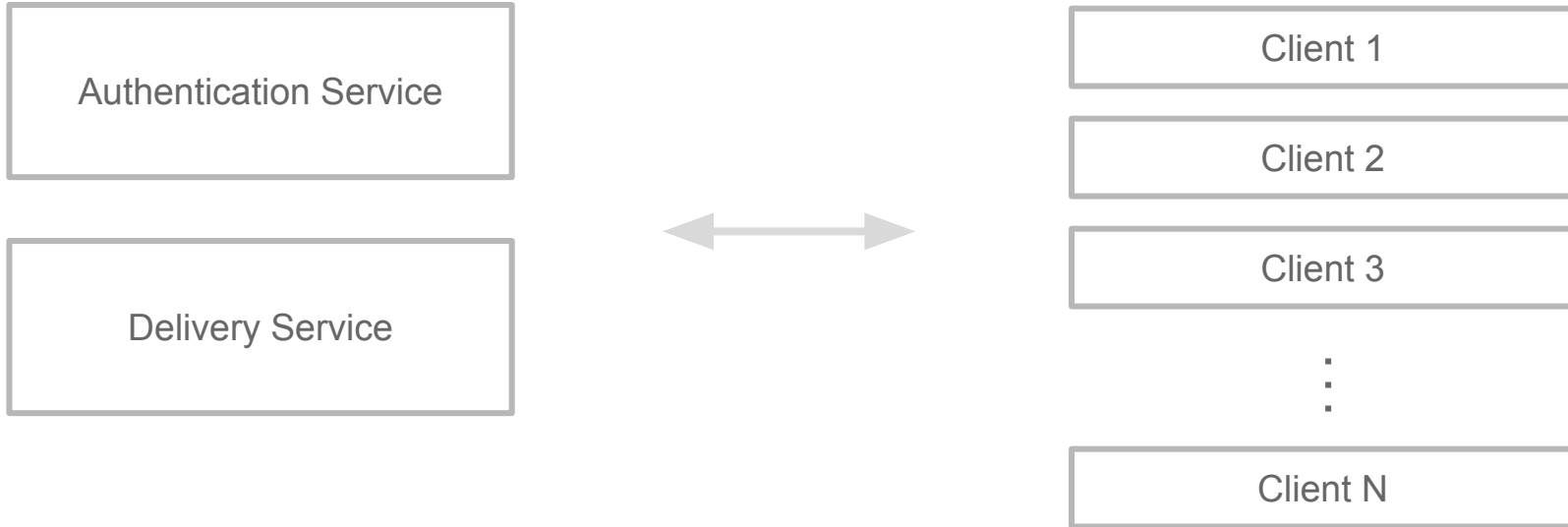


Forward Secrecy & Post-Compromise Security

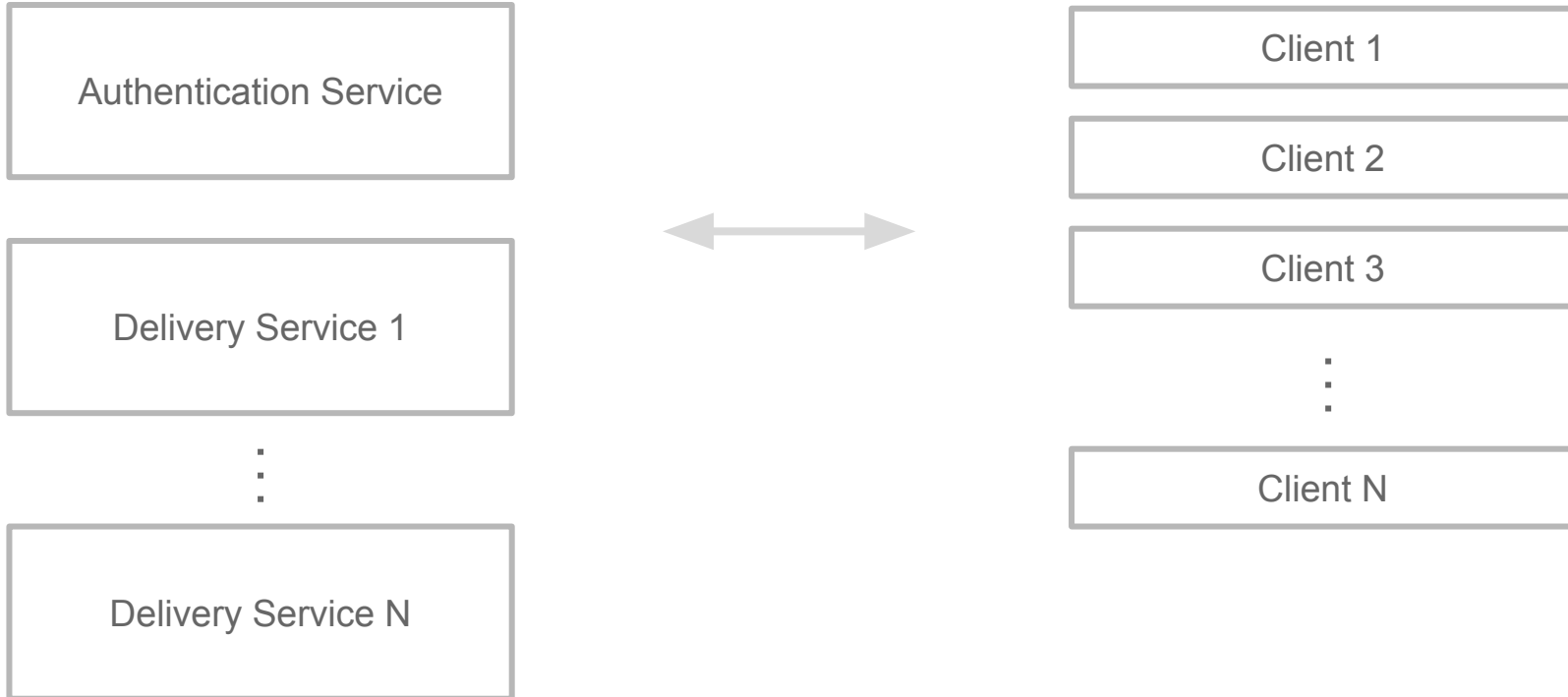




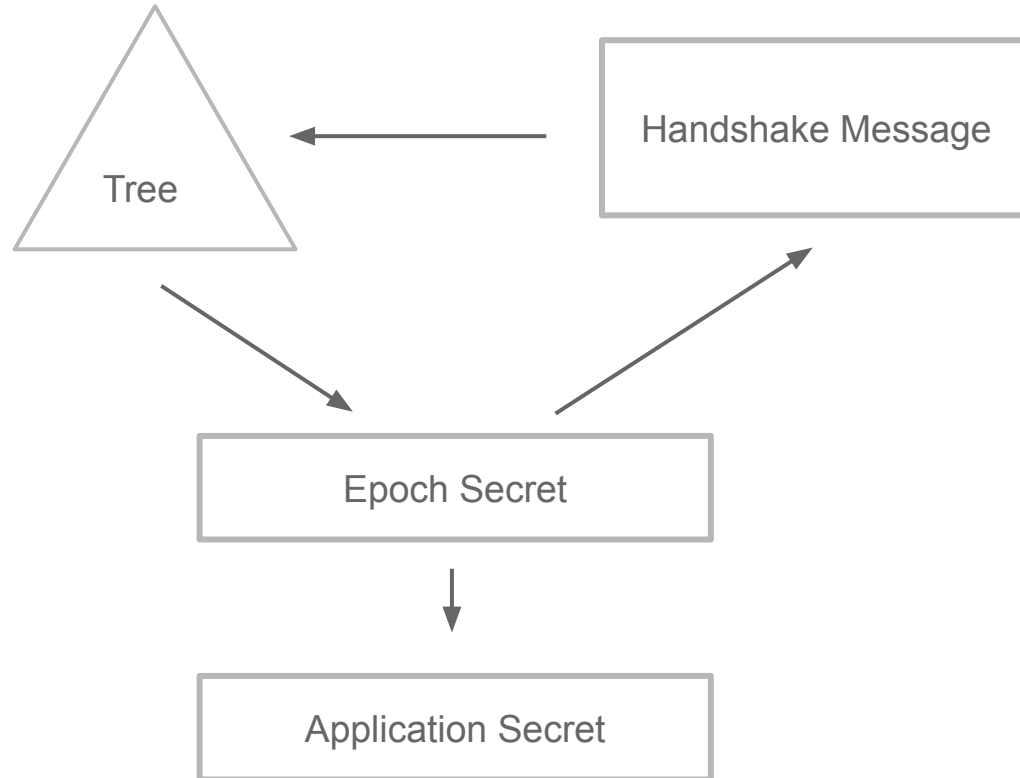
Architecture



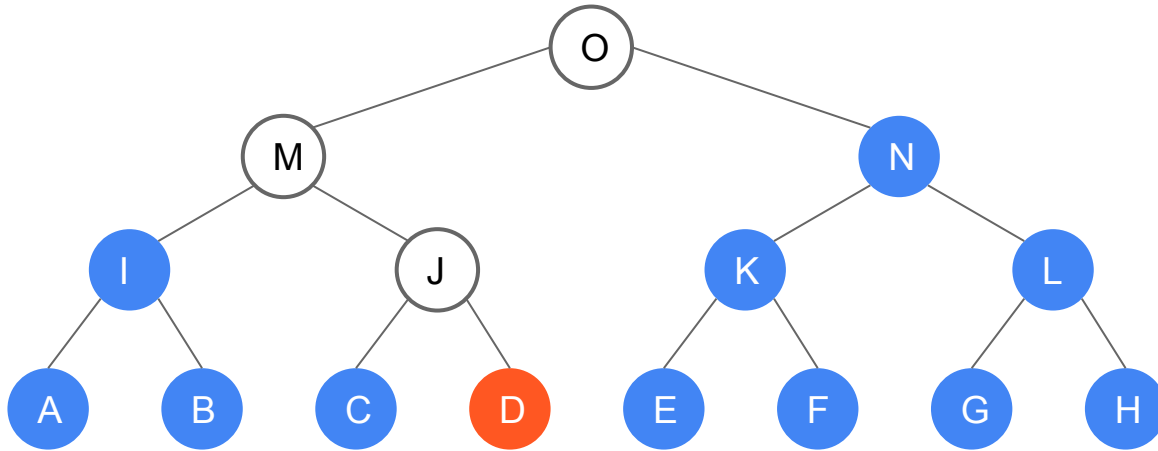
Federation



Protocol

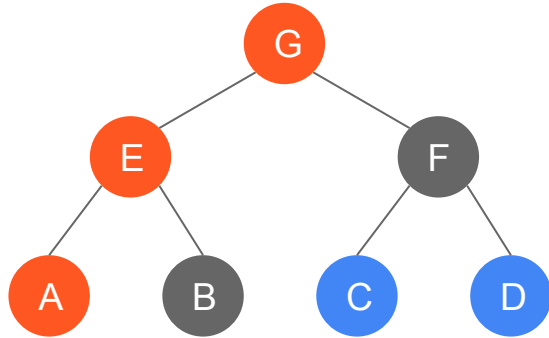


Tree invariant



The private key for an intermediate node is known to a member iff the node is an ancestor of the member's leaf

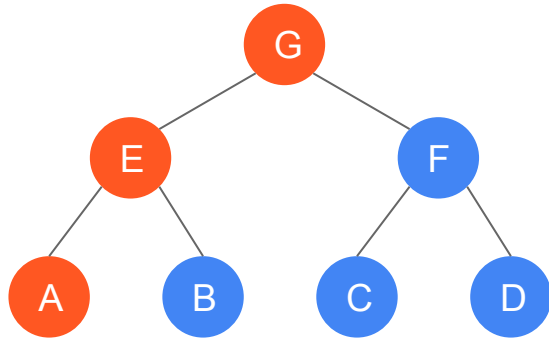
ART



$$E = \text{DH}(A, B)$$

$$G = \text{DH}(E, F)$$

TreeKEM



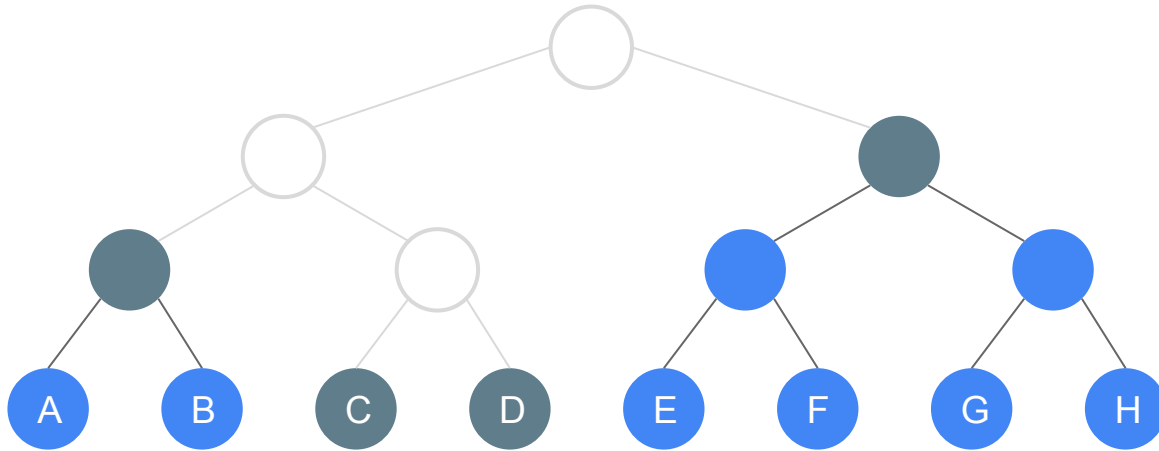
$$E = H(A)$$

$$G = H(E) = H(H(A))$$

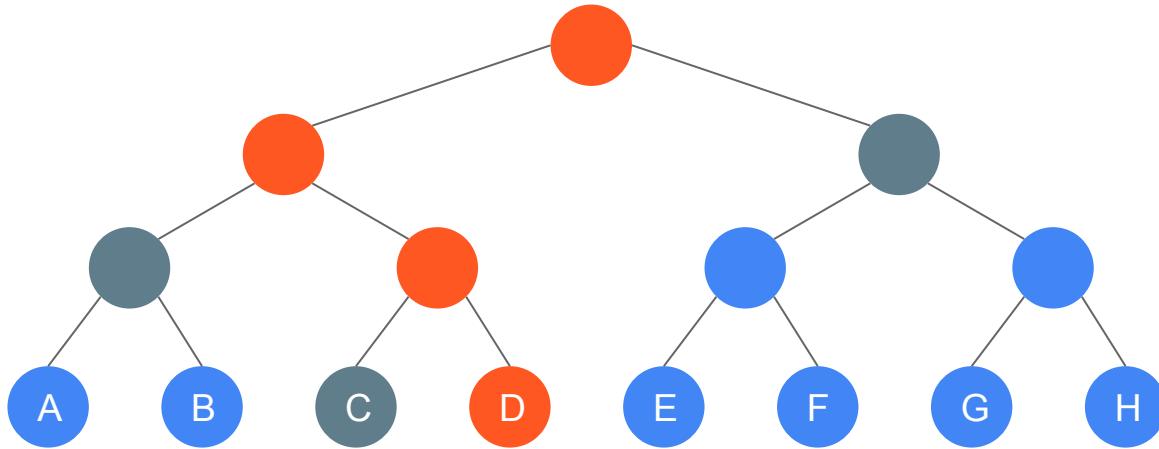
Advantages of TreeKEM

- Trees can contain blank nodes
- No “double joins”
- Updates can be processed in $O(1)$ instead of $O(\log N)$

Blank nodes



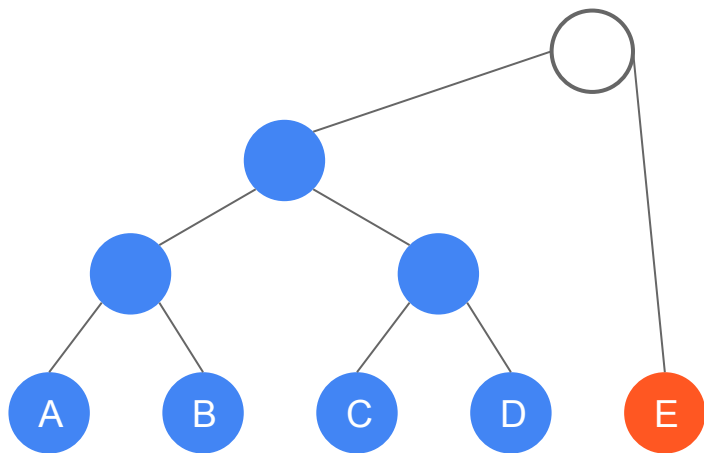
Update



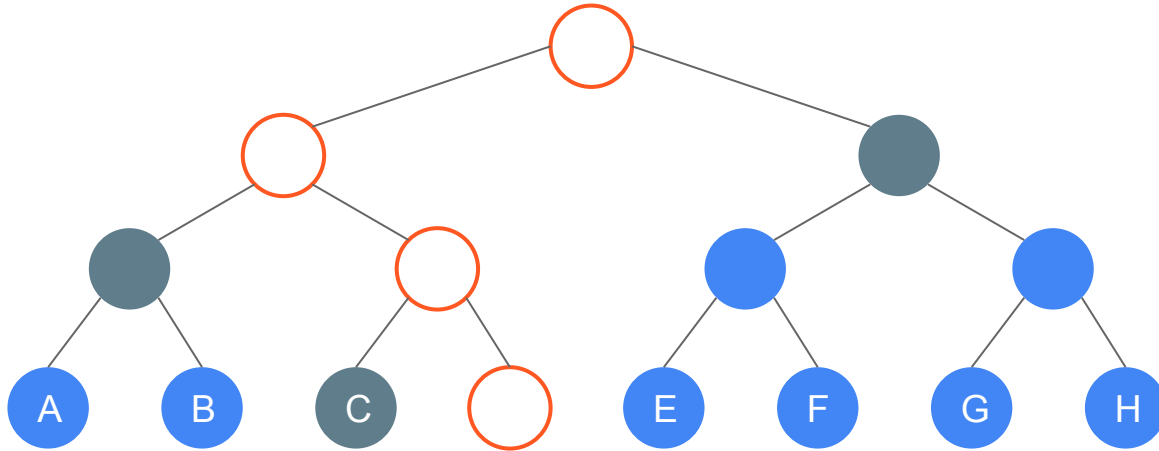
Direct path

Copath

Add



Remove



Direct path

Copath

Epochs

Alice creates group

Bob joins

Charlie joins

Bob leaves

Epoch 1
Secret 1

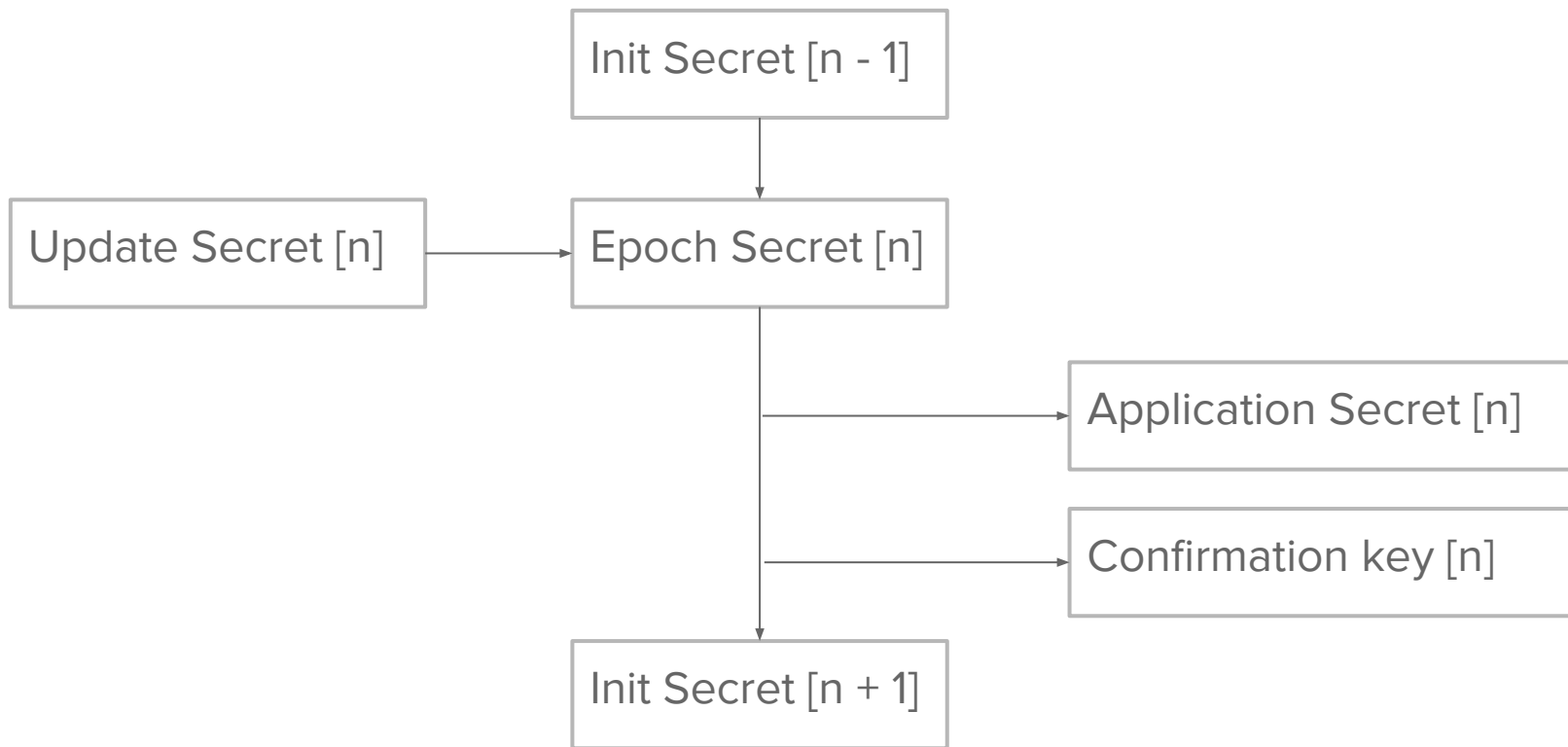
Epoch 2
Secret 2

Epoch 3
Secret 3

Epoch 4
Secret 4



Key Schedule



Authentication

Sign & MAC

- Handshake messages are signed with a signature key belonging to a client (keys are exchanged via Authentication Service)
- Handshake messages are authenticated with a MAC (confirmation key from the key schedule)

Working areas

- **Stabilizing** parts of the protocol for analysis
- **Security/Privacy:** HS message protection, general meta-data protection, malicious insider protection
- **Efficiency improvements:** Rebalancing trees, LazyUpdates, Tree warm-up, ...
- **Functionality improvements:** server initiated Add/Remove, recovery from state loss
- **Server:** Message ordering, ACK/NACK, federation, ...
- **Analysis:** TreeKEM & Authentication, general protocol, computational and symbolic verification, ...

Implementations

- **melissa** (Wire, Rust) <https://github.com/wireapp/melissa>
- **mlspp** (Cisco, C++) <https://github.com/cisco/mlspp>
- **MLS*** (Inria, F*)
- **molasses** (Trail of Bits, Rust) <https://github.com/trailofbits/molasses>
- **RefMLS** (NYU Paris, JS)
-  (Google, C++)

Messaging Layer Security

- Architecture: <https://github.com/mlswg/mls-architecture>
<https://architecture.messaginglayersecurity.rocks>
- Protocol: <https://github.com/mlswg/mls-protocol>
<https://protocol.messaginglayersecurity.rocks>
- Code + Interop: <https://github.com/mlswg/mls-implementations>
- Discussion: mls@ietf.org (archives)

https://app.wire.com/join/?key=qmrRRfakIMRm8UsYSqpA&code=KD8O6_Pvkli3pmzXbWtr

Thanks

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